

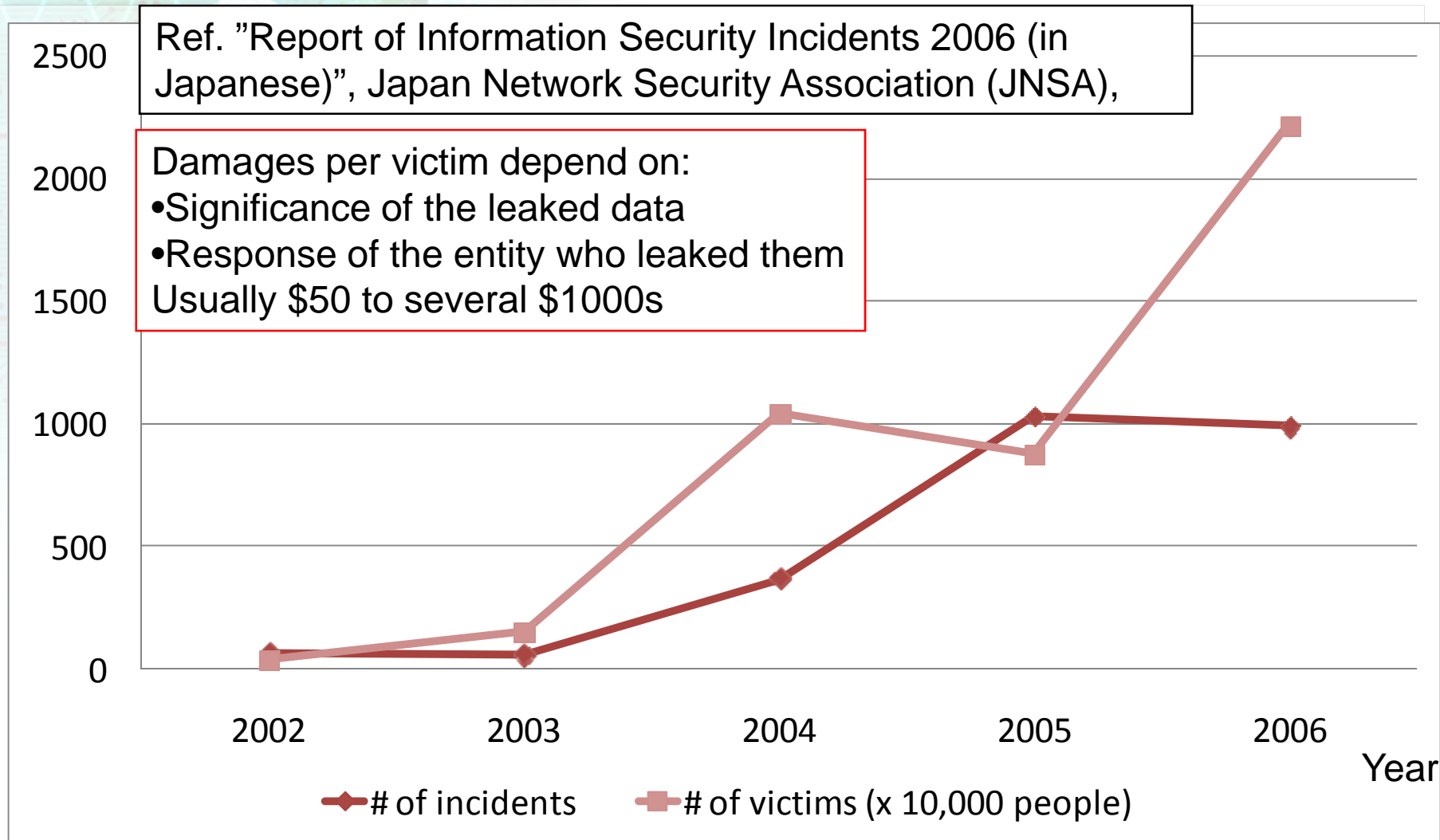
How to Cope with Information Leakage in The Ubiquitous Environment

Kaz Kobara^{1,2} and Hideki Imai^{2,1}

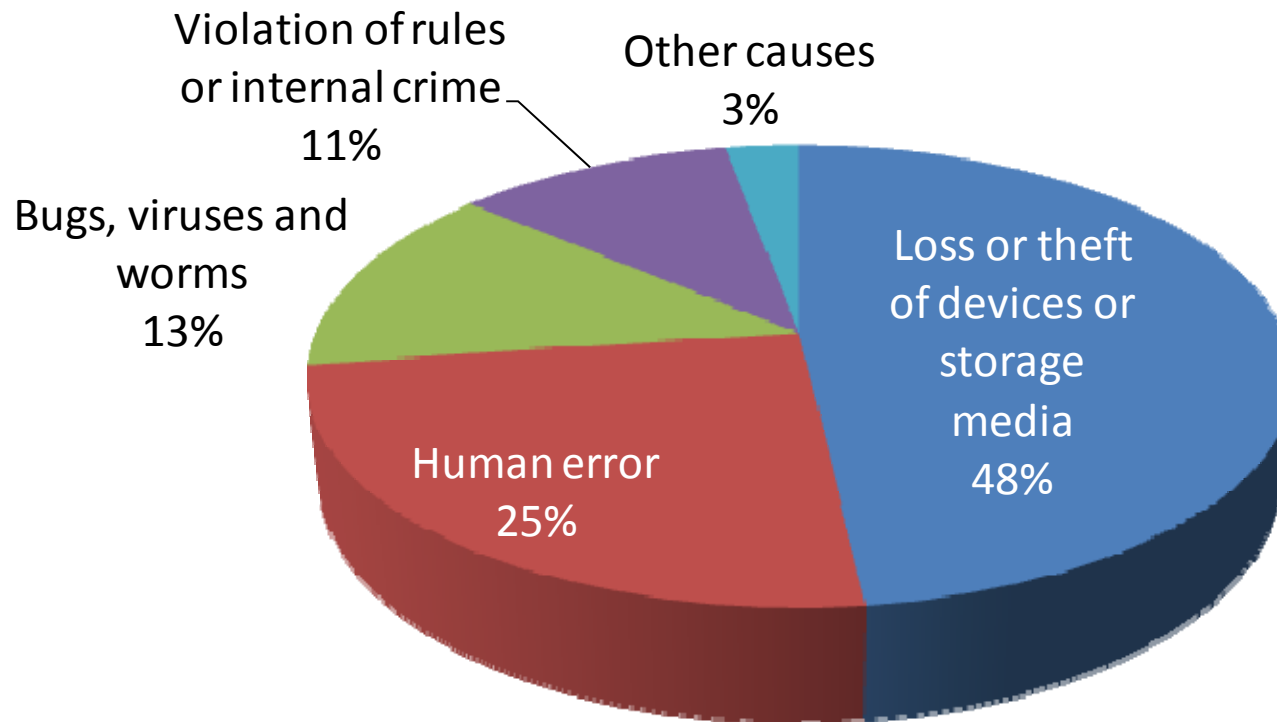
1: Research Center for Information Security (RCIS),
Advanced Industrial Science and Technology (AIST)

2: Chuo University

of Personal Data Leakage Incidents and Victims in Japan

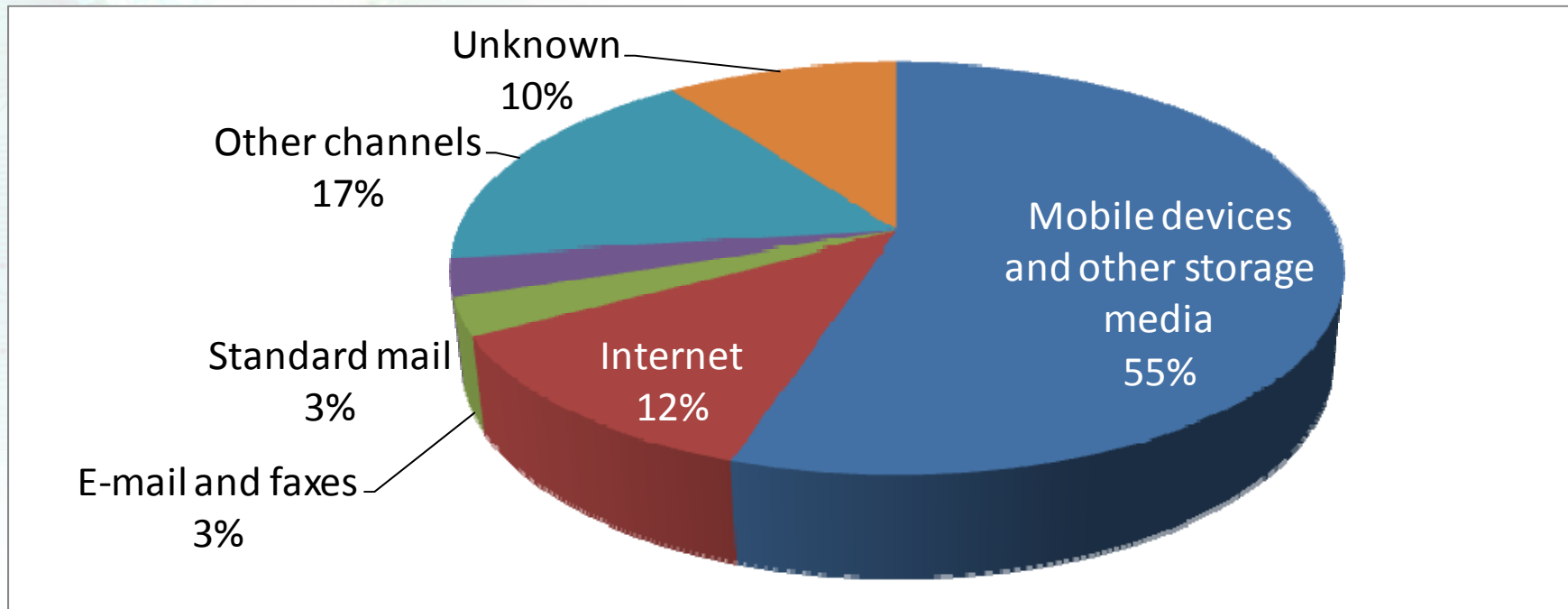


Causes of Leaks (Japan 2006)



Ref. "Report of Information Security Incidents 2006 (in Japanese)", Japan Network Security Association (JNSA),

Channels of Leaks (Across the Globe 2006)



“Global Data Leakage Survey 2006,” InfoWatch,
<http://www.infowatch.com/threats?chapter=162971949&id=207784626>

- Protection of mobile devices and storage media is important to resist against information leakage

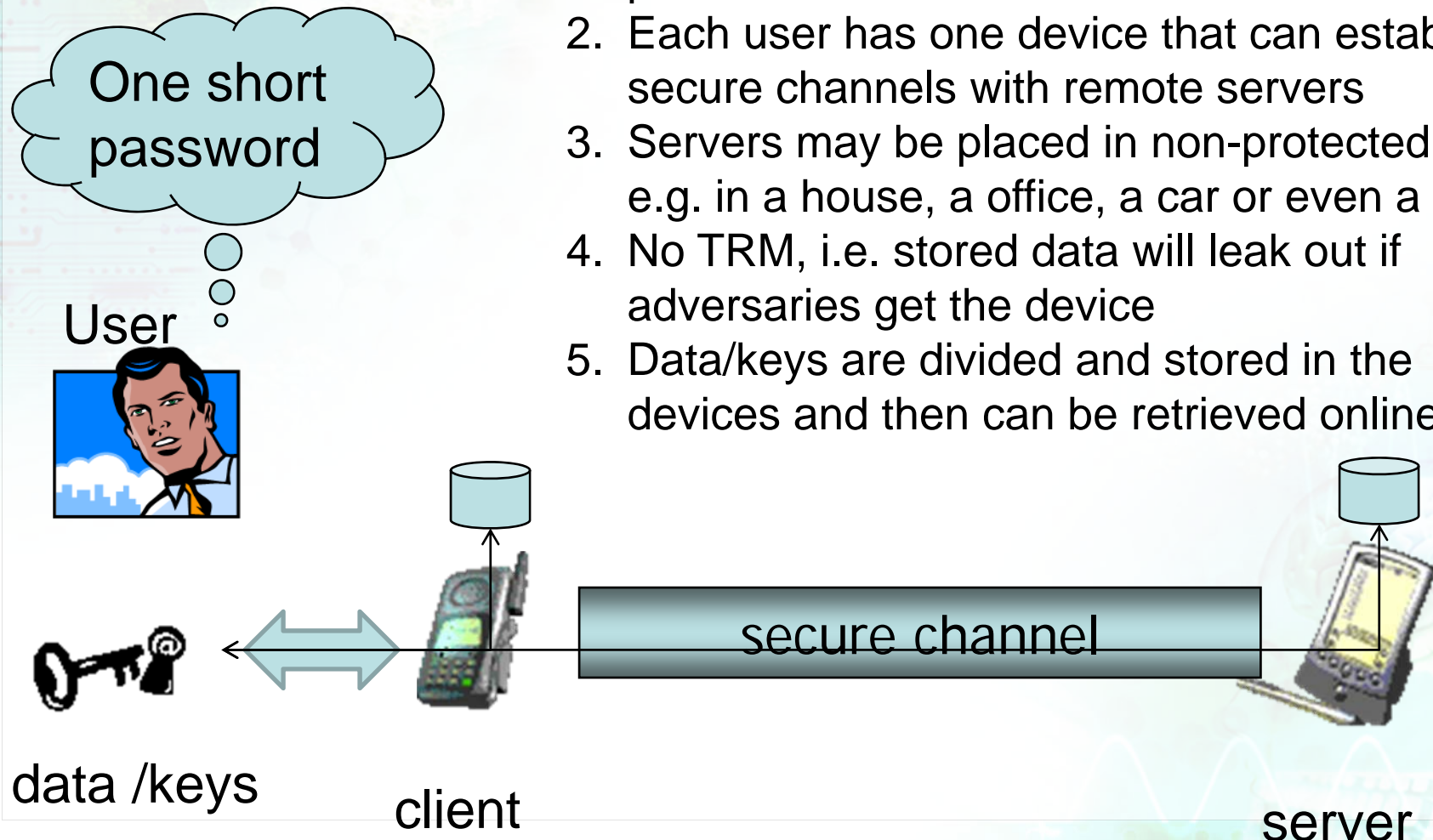
- Question is how to protect them?
 - One solution must be encryption but the **problem** is where to store the decryption key

- How about storing it in TRM?
 - It is still hard to realize perfect TRM with low-cost due to side channel attacks, such as DPA
- How about encrypting it with a password?
 - Short passwords can easily be exhaustively searched
 - Long passwords are hard to remember

TRM: Tamper Resistant Module
DPA: Differential Power Analysis

Scenario I: Two Node Construction (2NC)

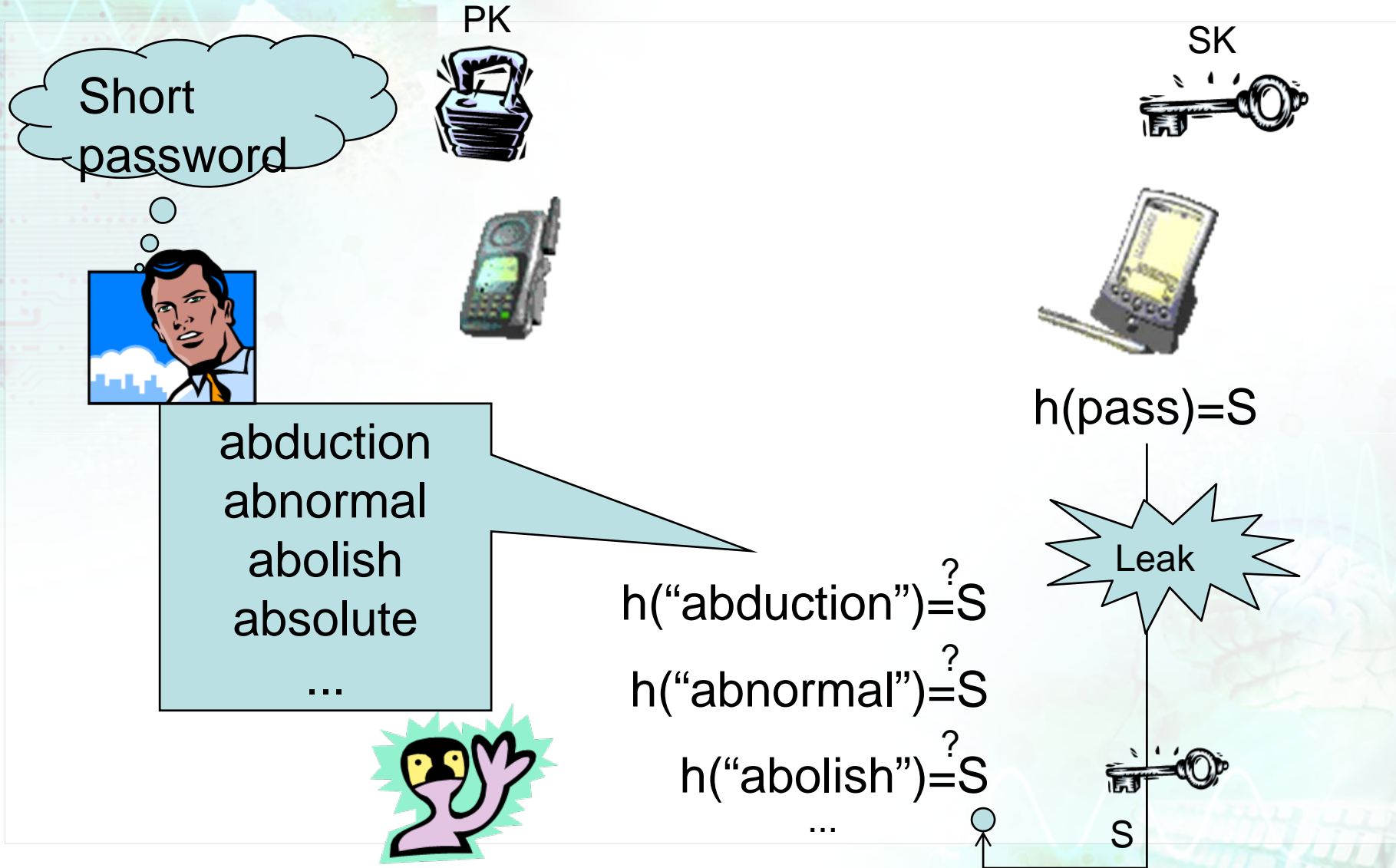
1. Each user remembers only one short password
2. Each user has one device that can establish secure channels with remote servers
3. Servers may be placed in non-protected area, e.g. in a house, a office, a car or even a bag
4. No TRM, i.e. stored data will leak out if adversaries get the device
5. Data/keys are divided and stored in the devices and then can be retrieved online



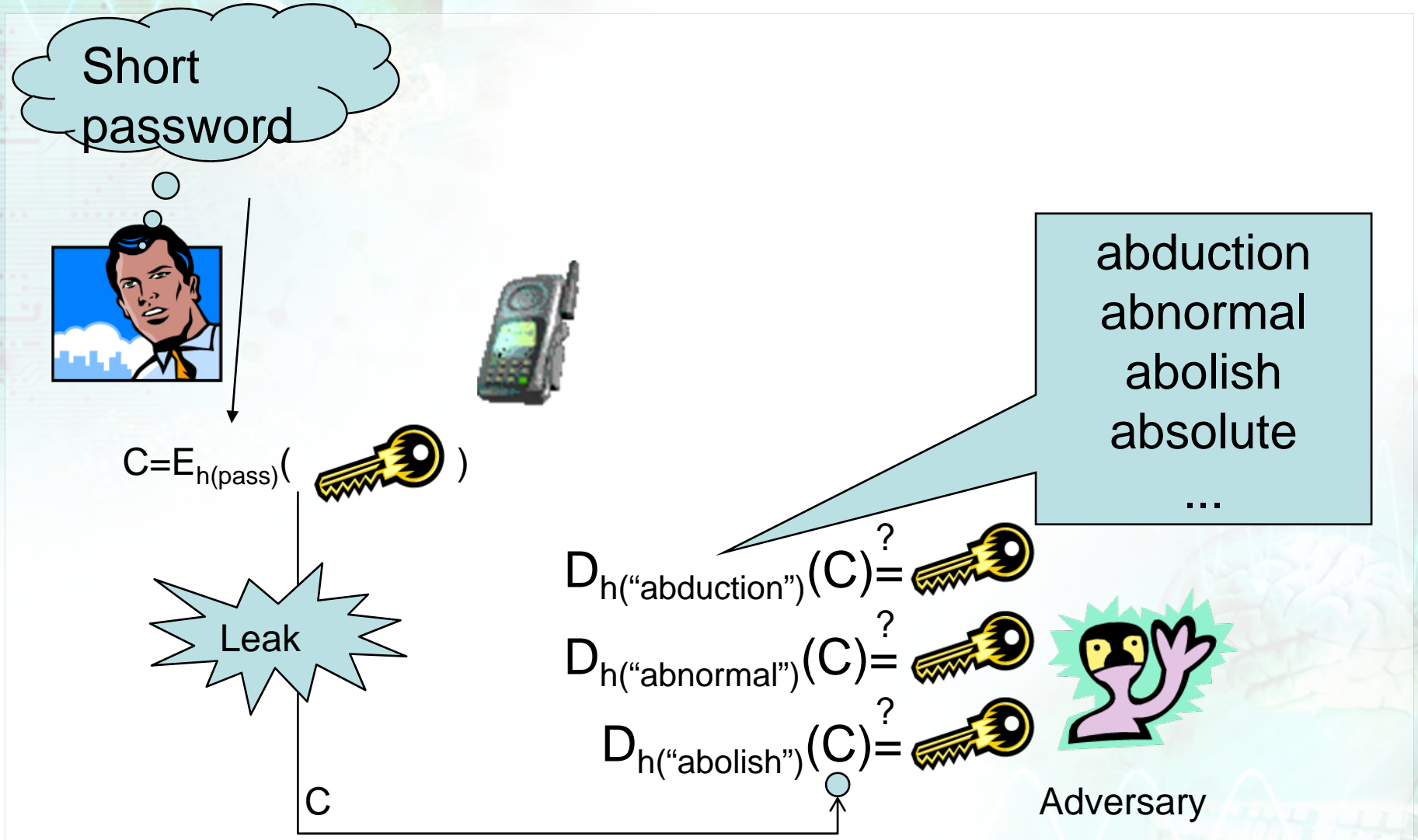
Problem 2

- While our scenario assumes leakage
- Most of the currently available protocols are vulnerable against information leakage
 - Since they are designed under the assumption that keys to establish secure channels are protected securely and never leak out
- So, once the keys leak out In their protocols,
 - adversaries can obtain the stored data or/and the user's personal password

Bad Example I (Hashed Password)



Bad Example II (PW-Protected-Keys)



Comparison among AKE Protocols

Can adversary obtain data or PW ? ○: No, X: Yes	Eaves dropping	Parallel On-line Attack	Resilience against Leakage			PW to remember
			From client	From server	From both with time difference	
Protocols						
Conventional PW-Only	X	X	○	X	X	Many
PAKE	○	X	○	X	X	Many
PKI (Server Auth.+PW)	○	X	○	X	X	Many
PKI (Server Auth.+PW+OTP)	○	○	○	X	X	Many
KPS+PW	○	○	X	X	X	One
PKI (Mutual Auth.)	○	○	X	○	X	One
LR-AKE (Our proposal)	○	○	○	○	○	One

LR-AKE (Leakage-Resilient AKE)

- New class of AKE (Authenticated Key Establishment) protocols
 - designed under the assumption that
 - Keys (more generally stored secrets) may leak out
 - can resist against information-leakage
- They fit with the scenario we consider

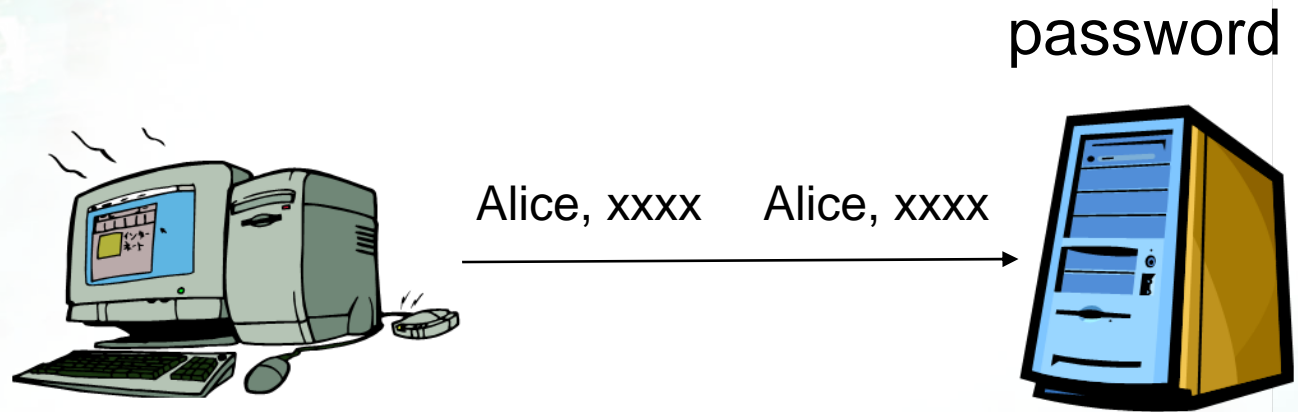
[SKI03] S. H. Shin, K. Kobara, and H. Imai, "Leakage-resilient authenticated key establishment protocols," Proc. of ASIACRYPT 2003, LNCS 2894, pp.166-172, 2003

[SKI07] S. H. Shin, K. Kobara, and H. Imai, "An Efficient and Leakage-Resilient RSA-Based Authenticated Key Exchange Protocol with Tight Security Reduction", IEICE Trans. Vol. E90-A, No. 2, pp. 474-490, 2007

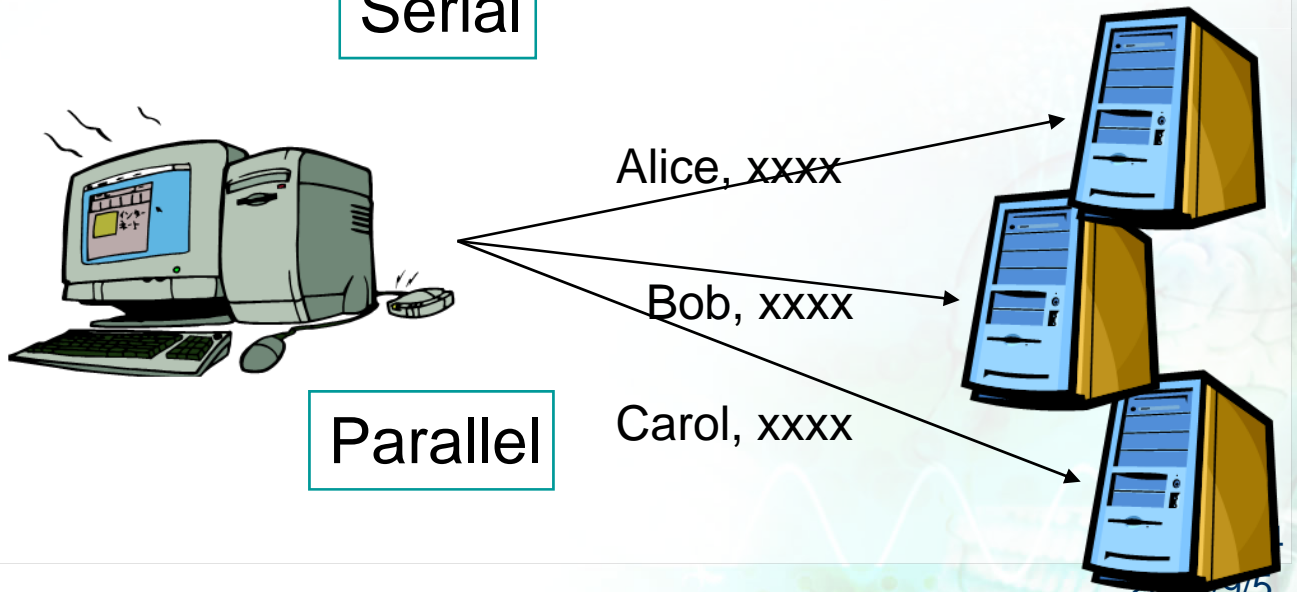
[NGSP] "New Generation Security Project," Ministry of Economy, Trade and Industry, 2005-2007

On-line Exhaustive Search

abduction
abnormal
abolish
absolute
...



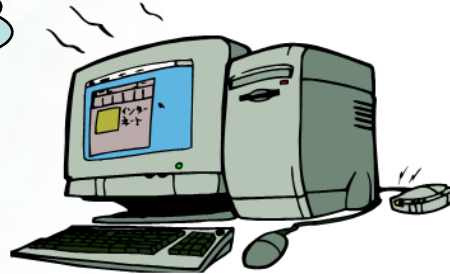
Serial



Parallel

Off-line Exhaustive Search

password



password



1. Gets the data for verifying the password
2. Tries password candidates off-line

This can be done with high speed in parallel

ID	PW
Alice	xxx
Bob	yyy

abduction
abnormal
abolish
absolute
...



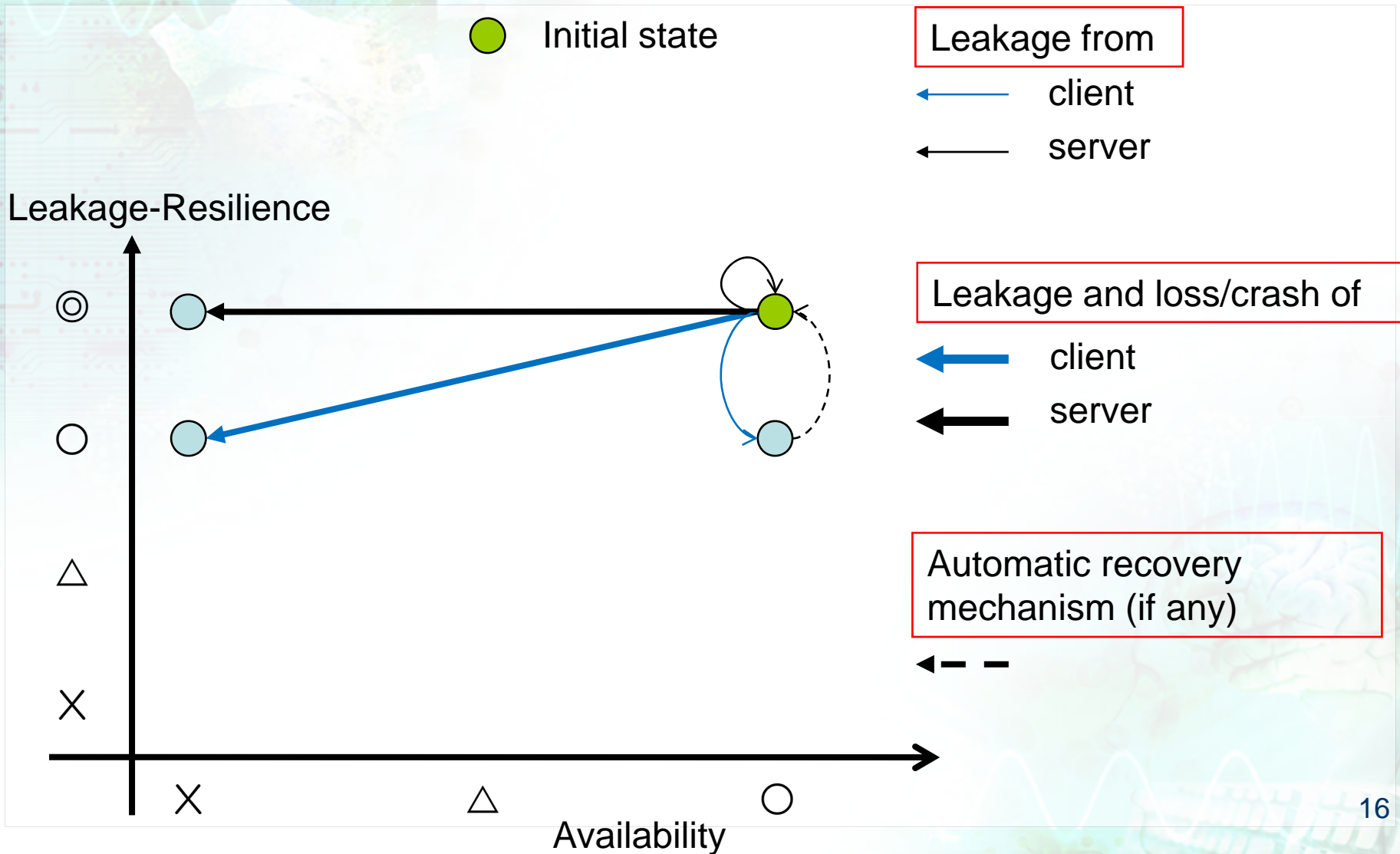
f("abduction") $\stackrel{?}{=} \text{XXX}$

f("abnormal") $\stackrel{?}{=} \text{XXX}$ XXX

f("abolish") $\stackrel{?}{=} \text{XXX}$

...

Damage against Node Compromise



Conventional PW-Only Protocols in 2NC

Such as CHAP, IPsec/IKE
(PSK), EAP-PSK and so on

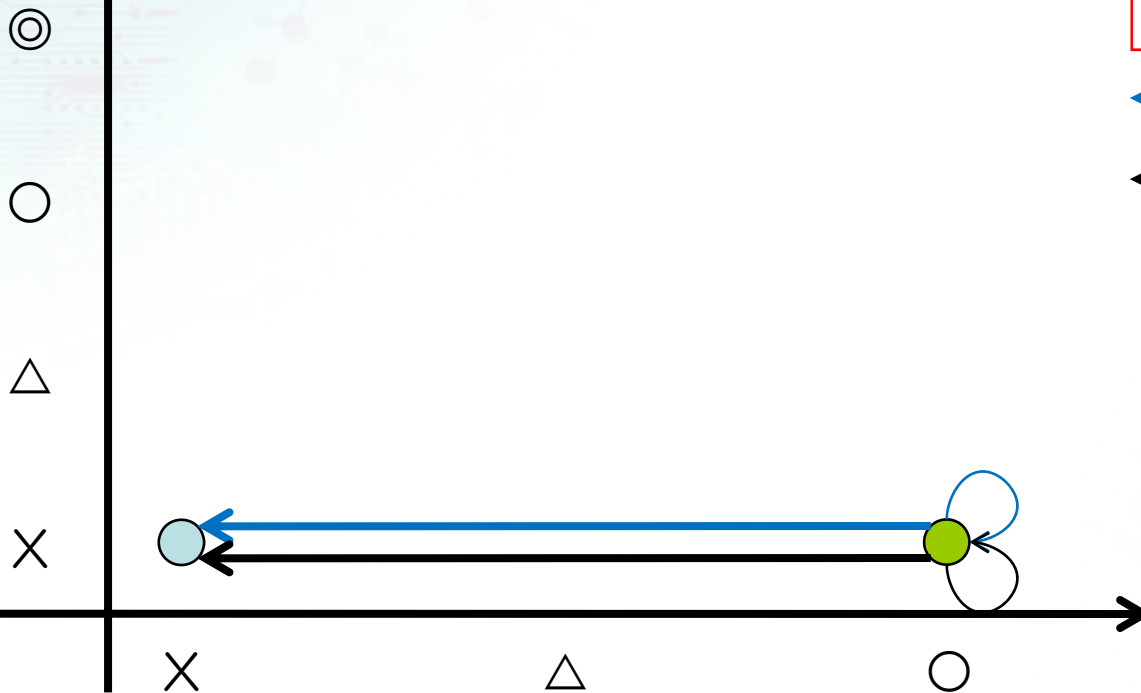
● Initial state

Leakage from

← client

← server

Leakage-Resilience



Leakage and loss/crash of

← client

← server

Availability

PKI (Server Auth+PW) or PAKE in 2NC

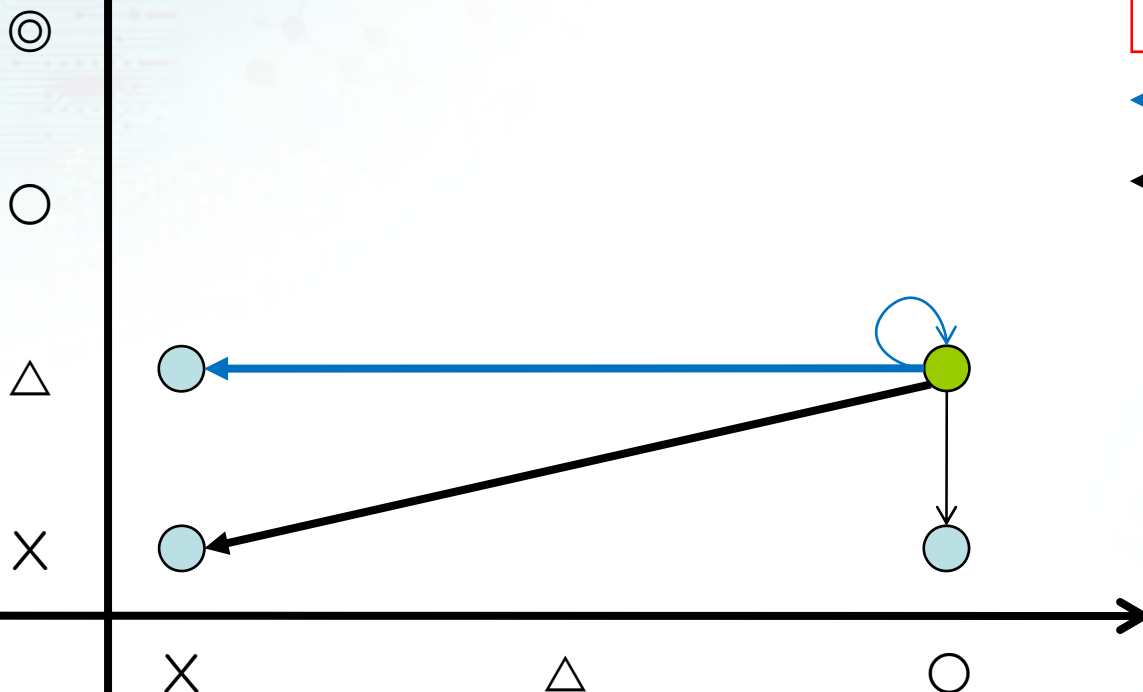
● Initial state

Leakage from

← client

← server

Leakage-Resilience



Leakage and loss/crash of

← client

← server

Availability

PKI (Server Auth+PW+OTP) in 2NC

● Initial state

Leakage from

← client

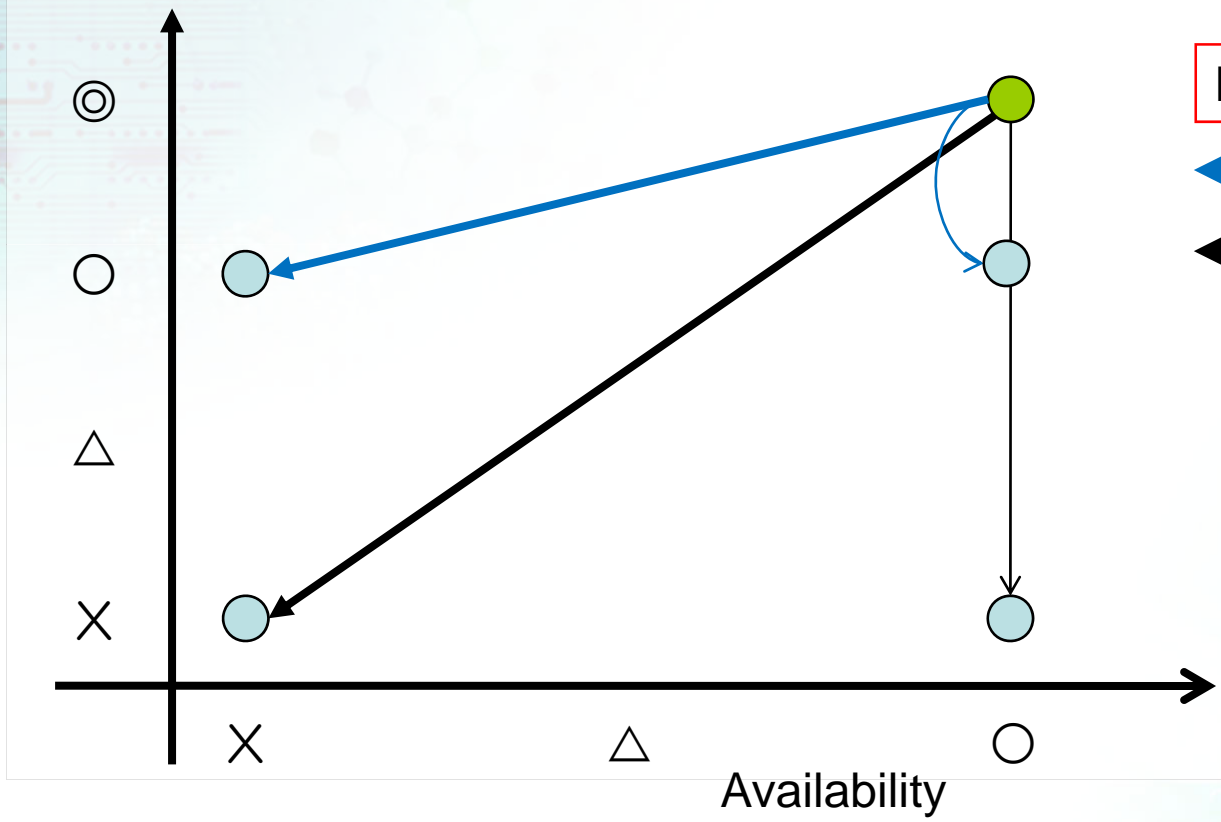
← server

Leakage-Resilience

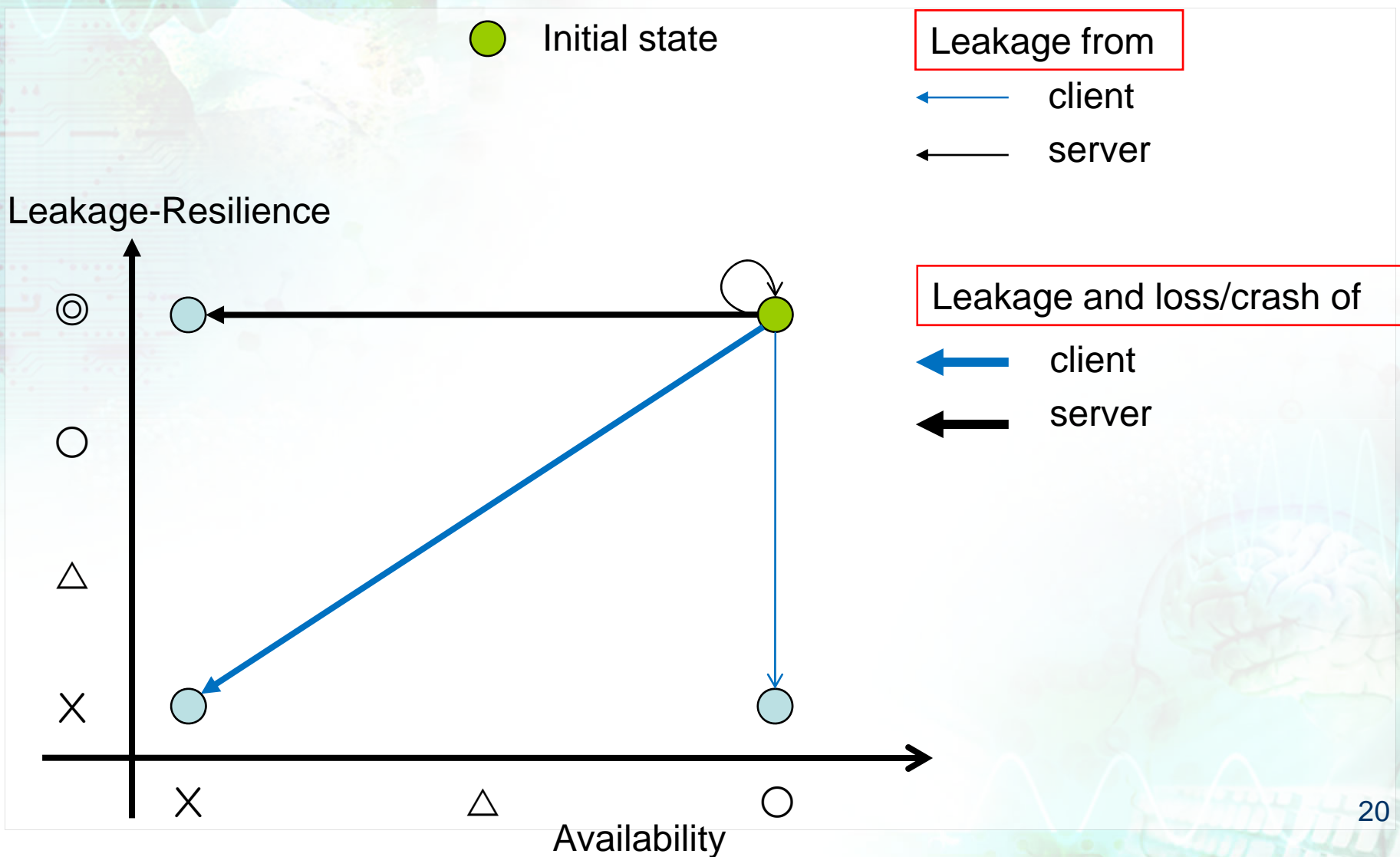
Leakage and loss/crash of

← client

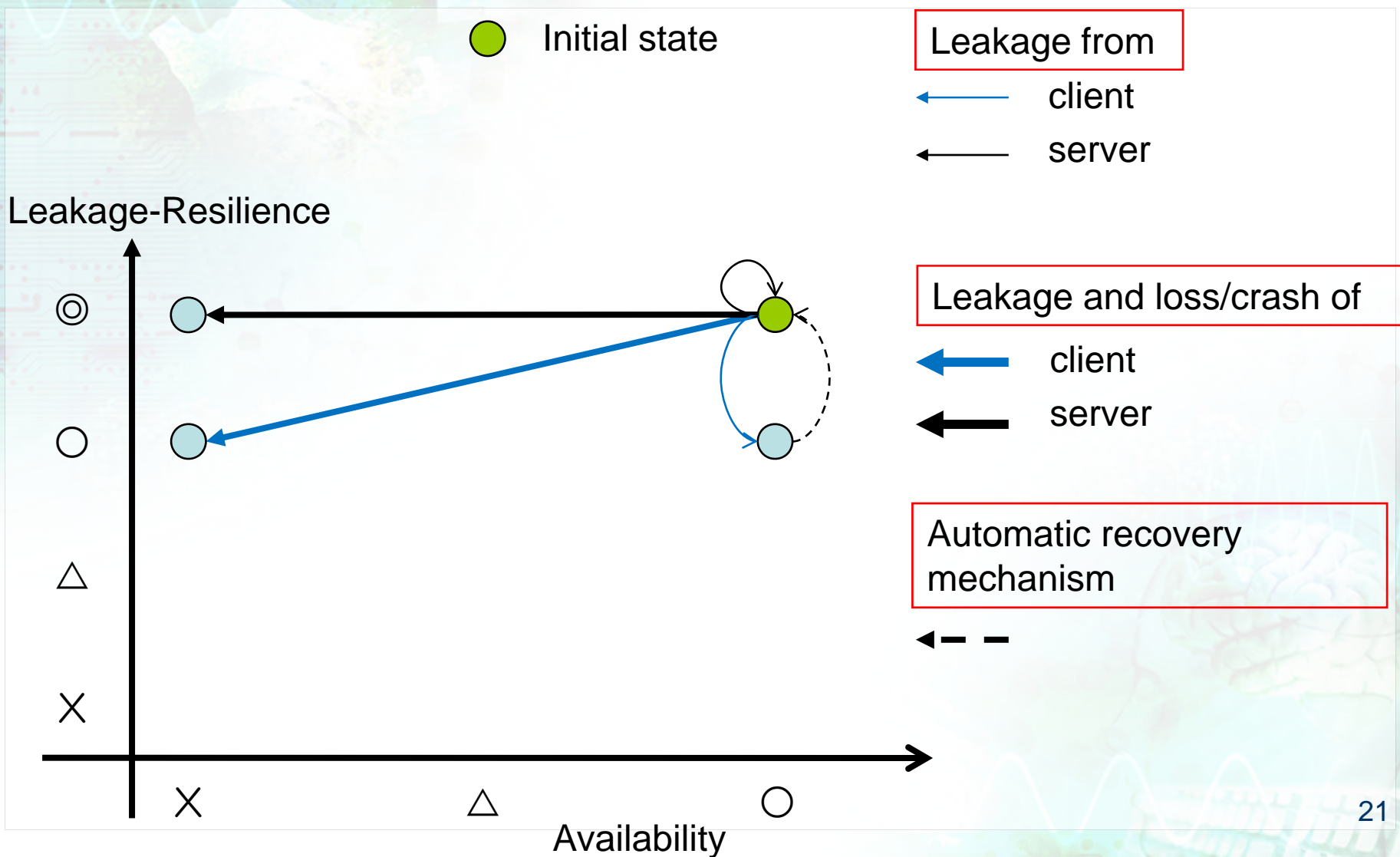
← server



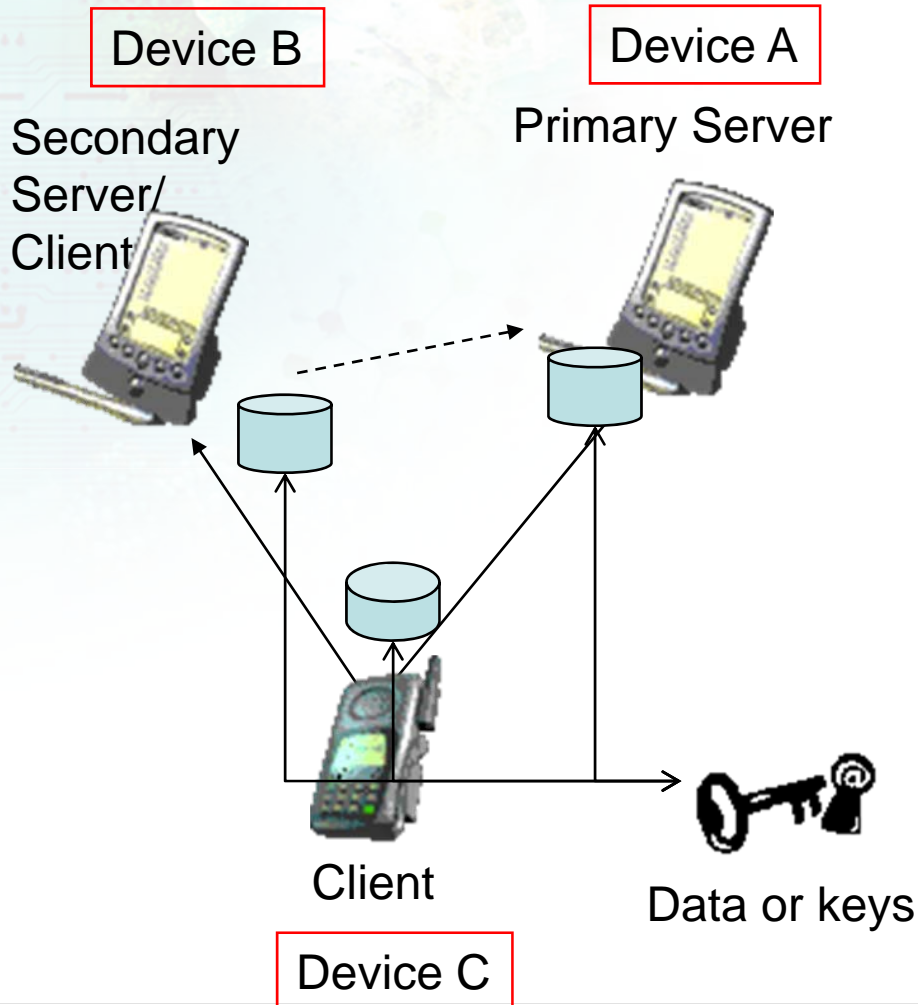
PKI (Mutual Auth) in 2NC



LR-AKE (Single Mode) in 2NC



Scenario II : Three Node Construction (3NC [Type A])



1. A user uses Device C as a client and Devices A and B as primary and secondary servers, respectively
2. When he/she lost Device C, visits at Device B and uses it as a client
3. Data/keys are divided and stored in these devices

PKI (Server Auth + PW) in 3NC

● Initial state

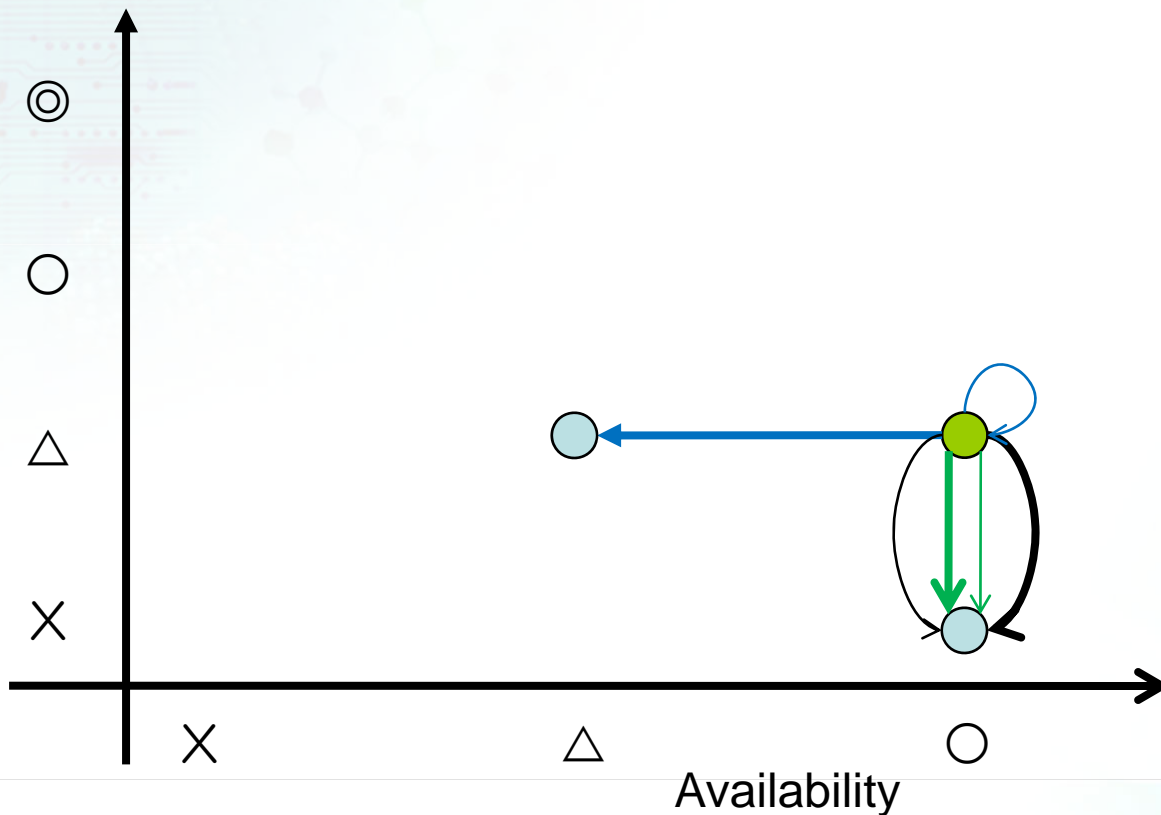
Leakage from

- ← client
- ← primary server
- ← secondary server

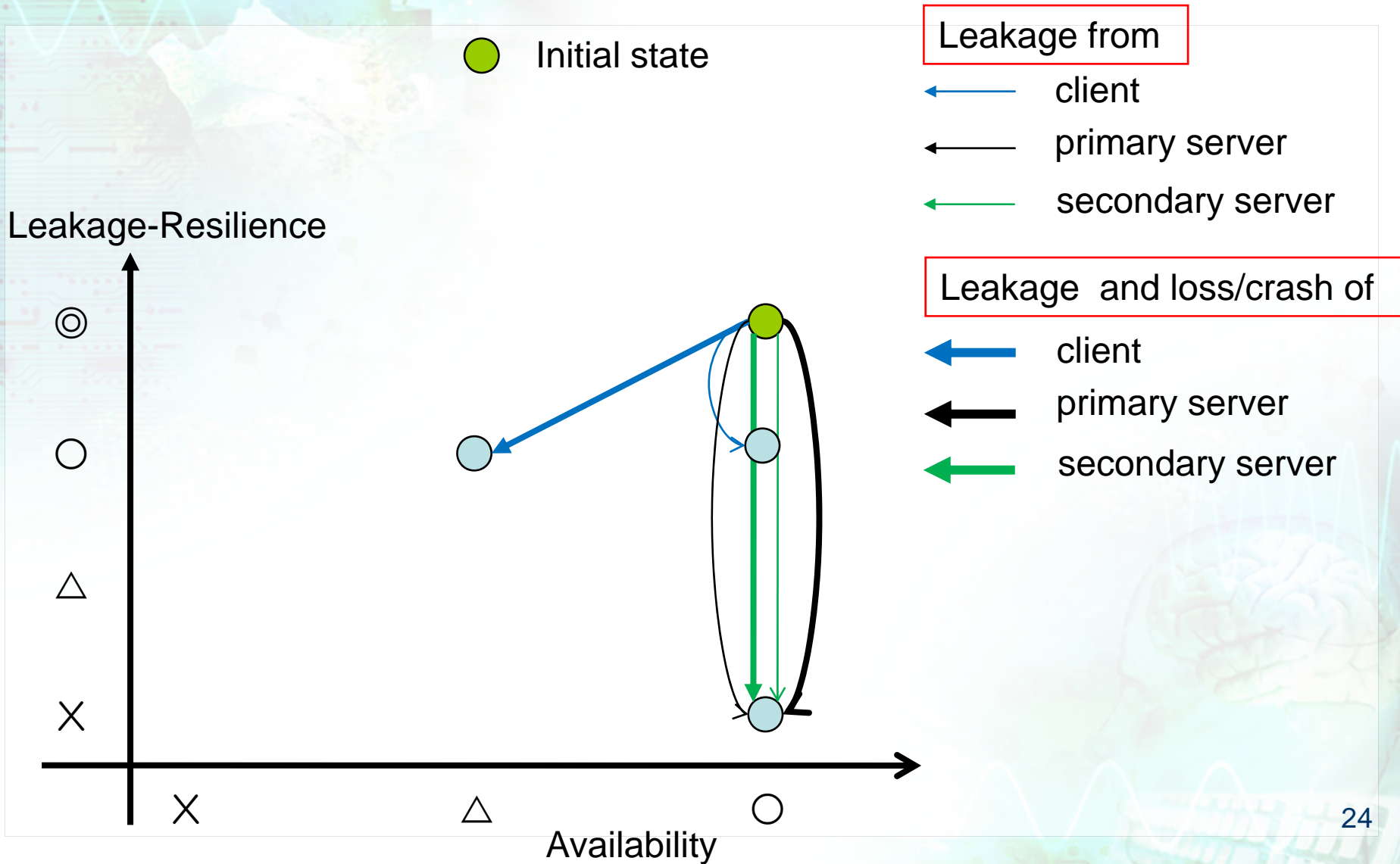
Leakage and loss/crash of

- ← client
- ← primary server
- ← secondary server

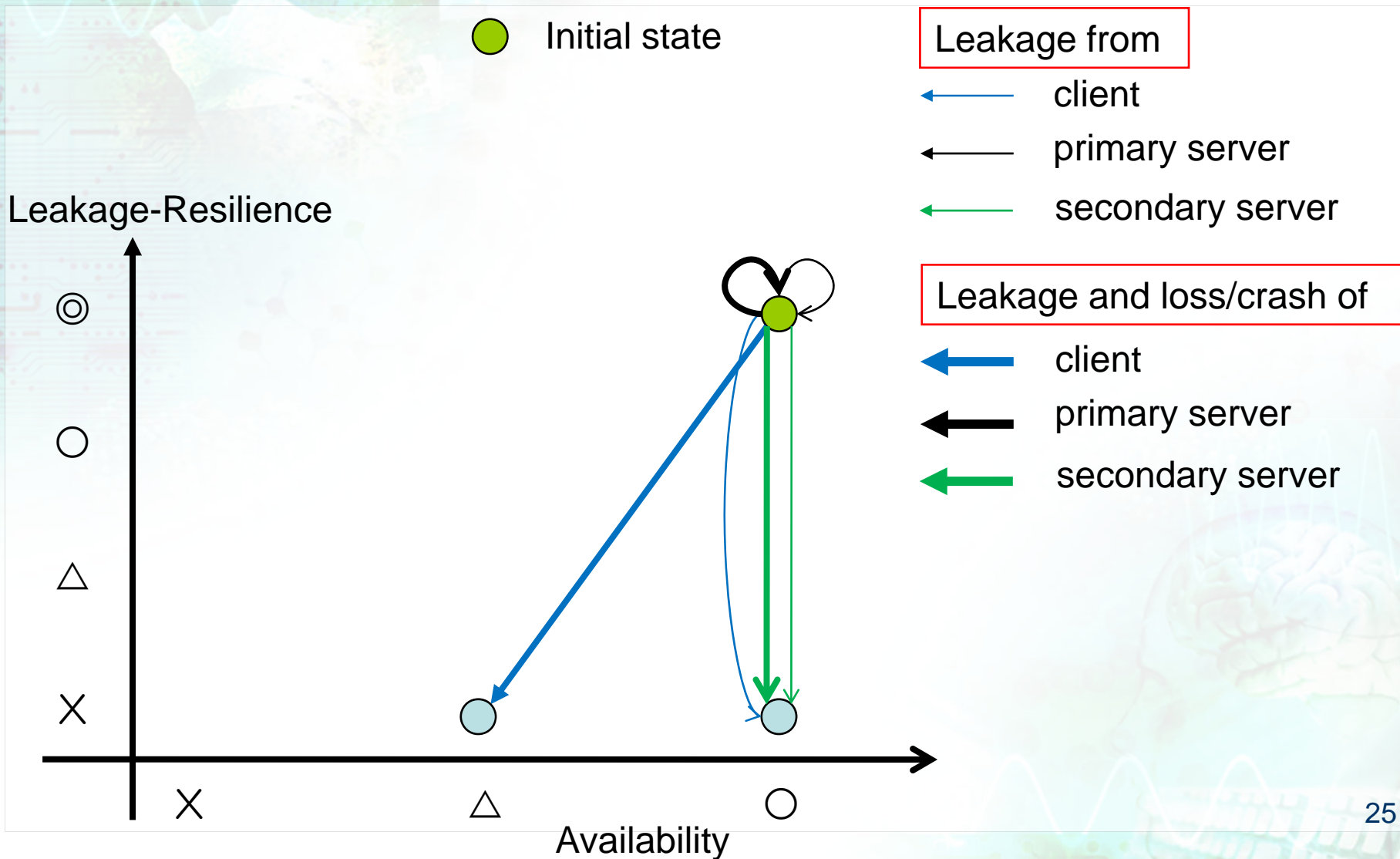
Leakage-Resilience



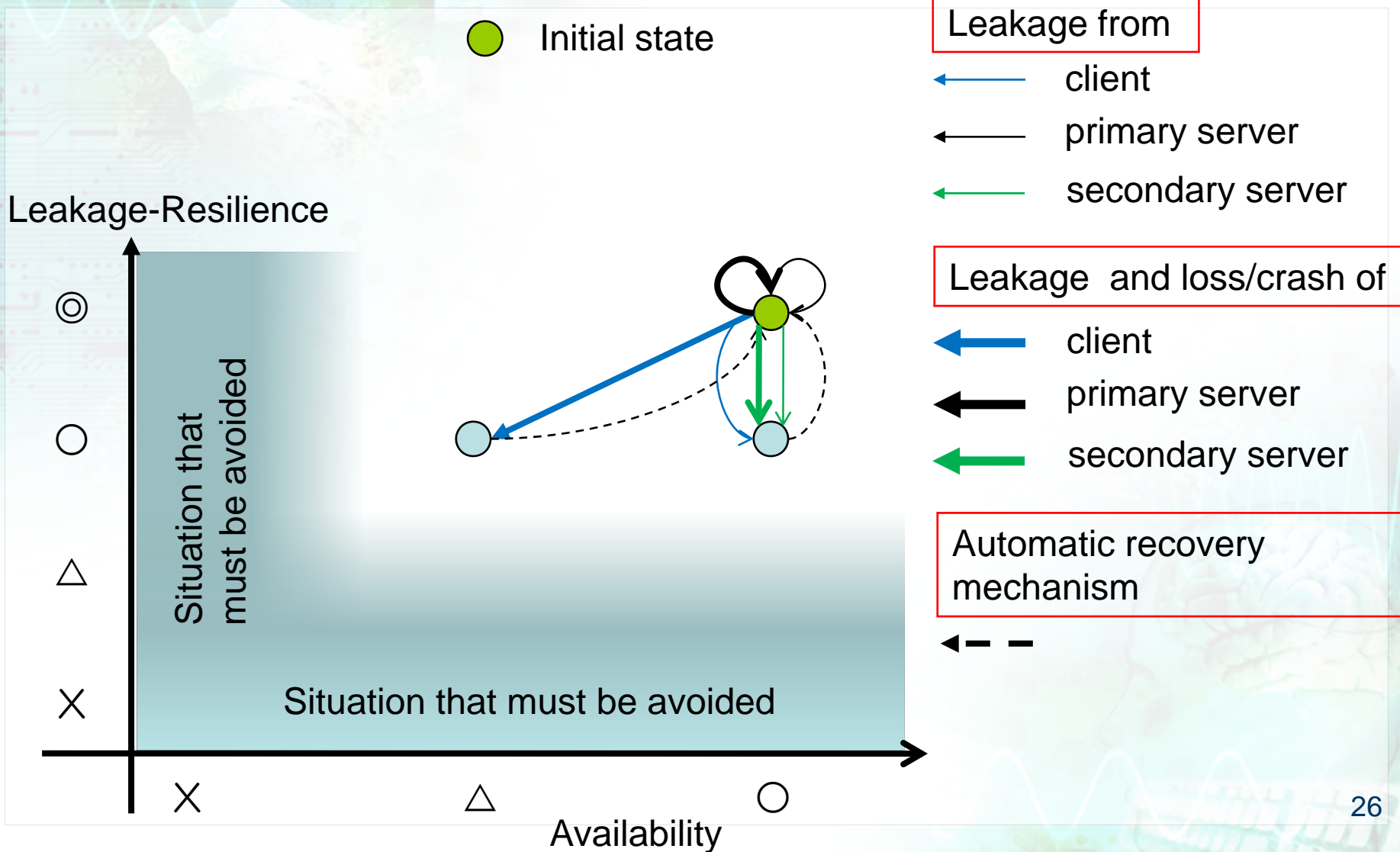
PKI (Server Auth+PW+OTP) in 2NC



PKI (Mutual-Auth) in 3NC



LR-AKE (Cluster Mode) in 3NC



- Leakage of critical information causes serious problems
- Encryption may be a solution, but the problem is where to store the decryption key
- We considered to store it in a distributed network
- And then showed the relationship of leakage resilience and availability
 - 3NC using LR-AKE has the best leakage resilience and availability